

# Announcements

1) Math Advising Session

11:30 - 1:30 CB 2047

(Math Library) - There  
will be pizza!

2) No online class

Back to Theorem:

Let  $T \in \mathcal{T}^m(\mathbb{R}^n)$ ,  $S \in \mathcal{T}^k(\mathbb{R}^n)$ ,

$R \in \mathcal{T}^l(\mathbb{R}^n)$ . Then

1) If  $A_{i+}(S) = 0$ , then

$$A_{i+}(T \otimes S) = 0.$$

2)  $A_{i+}(A_{i+}(T \otimes S) \otimes R)$

$$= A_{i+}(T \otimes A_{i+}(S \otimes R))$$

$$= A_{i+}(T \otimes S \otimes R) \quad (\text{associativity of } A_{i+})$$

3) " $\wedge$ " is associative, and

$$(\tau \wedge s) \wedge r$$

$$= \tau \wedge (s \wedge r)$$

$$= \frac{(k+m+l)!}{k!m!l!} \text{Alt}(\tau \otimes s \otimes r)$$

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1) is on last class's notes - involves some group theory. Assuming this, let's prove 2) and 3).

2) Consider

$$\overline{T \otimes S} - \text{Alt}(T \otimes S) = Q$$

Then

$$\text{Alt}(Q)$$

$$= \text{Alt}(T \otimes S) - \text{Alt}(\text{Alt}(T \otimes S))$$

(Alt linear)

$$= \text{Alt}(T \otimes S) - \text{Alt}(T \otimes S)$$

(Alt(Alt) = Alt)

$$= 0.$$

So by 1),

$$0 = \text{Alt}(Q \otimes R)$$

$$= \text{Alt}((T \otimes S - \text{Alt}(T \otimes S)) \otimes R)$$

$$= \text{Alt}(T \otimes S \otimes R - \text{Alt}(T \otimes S) \otimes R)$$

$$= \text{Alt}(T \otimes S \otimes R) - \text{Alt}(\text{Alt}(T \otimes S) \otimes R)$$

$$\Rightarrow \text{Alt}(T \otimes S \otimes R)$$

$$= \text{Alt}(\text{Alt}(T \otimes S) \otimes R)$$

The equality

$$A \mid + (T \otimes A \mid + (S \otimes R))$$

$$= A \mid + (T \otimes S \otimes R)$$

is proved similarly.

$$3) (T \wedge S) \wedge R$$

$$= \left( \frac{(m+k)!}{m!k!} A \mid (T \otimes S) \right) \wedge R$$

$$= \frac{(m+k)!}{m!k!} (A \mid (T \otimes S) \wedge R)$$

$$= \frac{\cancel{(m+k)!} (m+k+l)!}{m!k! \cancel{(m+k)!} l!} A \mid (A \mid (T \otimes S) \otimes R)$$

$$= \frac{(m+k+l)!}{m!k!l!} A \mid (T \otimes S \otimes R)$$

by 2)

The other equality

$$T \wedge (S \wedge R)$$

$$= \frac{(m+k+l)!}{m!k!l!} \text{Alt}(T \otimes S \otimes R)$$

is proved similarly. □



Proposition: (permutations)

Recall  $\varphi_i: \mathbb{R}^n \rightarrow \mathbb{R}$  is a 1-form,

$$\varphi_i\left(\sum_{j=1}^n a_j e_j\right) = a_i, \quad 1 \leq i \leq n.$$

$$1) \quad \varphi_i \wedge \varphi_i = 0$$

$$\forall \quad 1 \leq i \leq n$$

2) If  $i \neq j$ ,

$$\varphi_i \wedge \varphi_j = -\varphi_j \wedge \varphi_i$$

$$1 \leq i, j \leq n$$

Proof: Both at once!

$$\varphi_i \wedge \varphi_j = 2 A1 + (\varphi_i \otimes \varphi_j).$$

$$(\varphi_i \wedge \varphi_j)(v_1, v_2)$$

$$= 2 A1 + (\varphi_i \otimes \varphi_j)(v_1, v_2)$$

$$= \sum_{\sigma \in S_2} \text{sign}(\sigma) \varphi_i(v_{\sigma(1)}) \varphi_j(v_{\sigma(2)})$$

$$= \varphi_i(v_1) \varphi_j(v_2) \text{ (identity)}$$

$$- \varphi_i(v_2) \varphi_j(v_1) \text{ (flip)}$$

Observe if  $i=j$ , we get

$$\begin{aligned} & \varphi_i(v_1)\varphi_i(v_2) - \varphi_i(v_2)\varphi_i(v_1) \\ & = 0. \end{aligned}$$

If  $i \neq j$ ,

$$\begin{aligned} & (\varphi_j \wedge \varphi_i)(v_1, v_2) \\ & = \varphi_j(v_1)\varphi_i(v_2) - \varphi_j(v_2)\varphi_i(v_1) \\ & = -(\varphi_i \wedge \varphi_j)(v_1, v_2). \quad \square \end{aligned}$$

Theorem.  $\wedge^k(\mathbb{R}^n)$

has dimension  $\binom{n}{k} = \frac{n!}{k!(n-k)!}$ .

A basis is given by

$$\{ \varphi_{i_1} \wedge \varphi_{i_2} \wedge \dots \wedge \varphi_{i_k} \}$$

where  $1 \leq i_j \leq n \forall j$  and

$$i_1 < i_2 < \dots < i_k.$$

Proof: Let  $\omega \in \Lambda^k(\mathbb{R}^n)$ .

We know  $\omega \in \mathcal{Z}^k(\mathbb{R}^n)$ .

Then we can write

$$\omega = \sum_{i_1, i_2, \dots, i_k} c_{i_1, i_2, \dots, i_k} \varphi_{i_1} \otimes \varphi_{i_2} \otimes \dots \otimes \varphi_{i_k}$$

$\forall 1 \leq i_j \leq n$  since we know

$\{\varphi_{i_1} \otimes \dots \otimes \varphi_{i_k}\}$  is a  
basis for  $\mathcal{Z}^k(\mathbb{R}^n)$

Since  $\omega \in \Lambda^k(\mathbb{R}^n)$ ,

$\text{Alt}(\omega) = \omega$ . Then

by linearity of  $\text{Alt}$ ,

$$\omega = \text{Alt}(\omega)$$

$$= \sum_{i_1, i_2, \dots, i_k} c_{i_1, i_2, \dots, i_k} \text{Alt}(\varphi_{i_1} \otimes \varphi_{i_2} \otimes \dots \otimes \varphi_{i_k})$$

Fix  $i_1, i_2, \dots, i_k \in \{1, 2, \dots, n\}$ .

(Consider  $\text{Alt}(\varphi_{i_1} \otimes \dots \otimes \varphi_{i_k})$ .)

Now by our theorem,

$$\text{Alt}(\varphi_{i_1} \otimes \dots \otimes \varphi_{i_k})$$

$$= \frac{1}{k!} (\varphi_{i_1} \wedge \dots \wedge \varphi_{i_k})$$

(Theorem allows us to drop all parentheses in wedge product) -

Case 1:  $\exists i_s, i_t, i_s = i_t$   
( $s \neq t$ ).

Using the second property of the proposition, we can produce

$\mathcal{Q}i_s \wedge \mathcal{P}i_t$  in our wedge product, which is then zero by the first property of the proposition since  $i_s = i_t$ .



This tells us that if

$$\varphi_{i_1} \wedge \dots \wedge \varphi_{i_k} \neq 0,$$

we must have

$$i_s \neq i_t \quad \text{for any} \\ 1 \leq s, t \leq k.$$

We have then reduced  
the dimension of  $\Lambda^k(\mathbb{R}^n)$

to no more than

$$n \cdot (n-1) \cdots (n-k) = P(n, k)$$

Case 2:  $i_s \neq i_t$  for any

$$1 \leq s, t \leq k.$$

Then by using property

2) from the proposition

repeatedly, we may

rearrange

$$\mathcal{C}_{i_1} \wedge \mathcal{C}_{i_2} \wedge \dots \wedge \mathcal{C}_{i_k} \text{ as}$$

$$\mathcal{C}_{j_1} \wedge \mathcal{C}_{j_2} \wedge \dots \wedge \mathcal{C}_{j_k}$$

Where

$$j_1 < j_2 < \dots < j_k$$

and we may have

to change the sign of  
the wedge product.

This says

$\varphi_{i_1} \wedge \dots \wedge \varphi_{i_k}$  is  
a scalar multiple of

$$\varphi_{j_1} \wedge \dots \wedge \varphi_{j_k}$$

This implies that

$$\dim(\wedge^k(\mathbb{R}^n)) \leq \binom{n}{k}.$$

To show linear independence

$$\text{of } \{ \varphi_{i_1} \wedge \varphi_{i_2} \wedge \dots \wedge \varphi_{i_k} \}$$

$$\text{with } i_1 < i_2 < \dots < i_k,$$

apply to basis vectors.

(exactly as in the proof  
for  $\varphi_{i_1} \otimes \varphi_{i_2} \otimes \dots \otimes \varphi_{i_k}$ )

Observe that

$$\{ \varphi_{i_1} \wedge \dots \wedge \varphi_{i_k} \}$$

$$i_1 < i_2 < \dots < i_k$$

has cardinality  $\binom{n}{k}$ .  $\square$

Orientation: From the

theorem just proved,

$$\dim(\Lambda^n(\mathbb{R}^n)) = \binom{n}{n} = 1.$$

Since  $\det \in \Lambda^n(\mathbb{R}^n)$ , any

$\omega \in \Lambda^n(\mathbb{R}^n)$  is a scalar

multiple of  $\det$ !

We could have  
defined  $\det$  as

the unique alternating  
 $n$ -tensor on  $\mathbb{R}^n$  satisfying

$$\det(e_1, e_2, \dots, e_n) = 1.$$

Any basis for  $\mathbb{R}^n$   
can be obtained by  
applying an invertible  
linear operator to  
 $\{e_1, \dots, e_n\}$ . So  
if  $\{f_1, \dots, f_n\}$  is such  
a basis and  $S$  is  
the linear operator,



$$\begin{aligned} \det(f_1, \dots, f_n) \\ &= \det(S) \det(e_1, \dots, e_n) \\ &= \det(S). \end{aligned}$$

So we break bases up  
into two categories:

$$\det(f_1, \dots, f_n) > 0 \text{ and}$$

$$\det(f_1, \dots, f_n) < 0.$$

An orientation on  $\mathbb{R}^n$  is a choice of either bases with positive determinant or base with negative determinant but not both!